

Branching strategies for solving pseudo-Boolean optimization problems using ILP solvers

(Pragmatics of SAT Workshop)

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Outline of the presentation

Preliminaries

Frequency based branching heuristic: MOHP

Heuristic for BCP based look-ahead branching

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- ▶ PBO: ILP with binary variables
 - Can be solved using ILP solvers such as CPLEX

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 - Poor scaling upon distribution [Ralphs, T., Shinano, Y., Berthold, T., Koch, T.(2018)]
 - Dependency on pseudo-cost updates from other parts of the search tree

Motivation and Goals

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 - Proposed heuristics can also be used in existing PBO solvers

Use case

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- ▶ ILP has many constraints with 2 variables (Boolean Constraint Propagation, BCP)

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 - Strong branching (ILP), but it is known that probing has a role to play in branching
- ▶ **First few branching decisions are vitally important ! (for both SAT and ILP)** [Heule, M. J., Kullmann, O., Wieringa, S., & Biere, A. (2011, December).]

Frequency based branching heuristic: MOHP

Maximal occurrence in Highest Priority

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- ▶ In the presence of an objective function, we change the priority definition that MOMS uses
 - quickly get rid of infeasible vertices which have very low objective function values

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 - also get rid of infeasible vertices in the vicinity that have low objective values

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- ▶ Branch on x_2 or x_3 , tie-break won by x_2

MOHP observations

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- ▶ In a best-first search, this vertex is also the optimal solution

Heuristic for BCP based look-ahead branching

Motivation

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- ▶ Example: PROP heuristic [Li and Anbulagan, 1997]

Dominated Trigger

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- ▶ Only variables from APEX triggers are considered for branching
- ▶ Branch on variables on the periphery of the implication graph

Apex triggers in the constraint set below are $(x_4, 0)$ and $(x_6, 1)$.

We branch on one of x_4 or x_6 , using maximin criteria

$$c1 : x_1 + x_2 \geq 1$$

$$c2 : x_3 + x_2 \leq 1$$

$$c3 : x_3 + x_5 \geq 1$$

$$c4 : x_7 + x_4 \geq 1$$

$$c5 : x_1 + x_7 \leq 1$$

$$c6 : x_5 + x_6 \leq 1$$

Table 1. A comparison of CPLEX strong branching with look-ahead based branching and MOHP.

PBO	Root node BCP			Results				Nodes (first hour)		Nodes (end of test)	
	Fractional	Apex	Savings	Time	Solution	Best-bound	Gap	Remaining	Processed	Remaining	Processed
LTL routing 11 th March	2	2	0	24	75863.82	73486.57	0.03	32180	36189	770272	940175
				24	75334.53	74550.1	0.01	767	1837	782504	1107662
				24	75196.82	74590.76	0.01	1137	1871	584775	1098021
LTL routing 9 th March	6	6	0	24	58391.92	54768.51	0.06	15631	20580	358302	400829
				24	58179.65	57230.78	0.02	4118	9580	161024	293392
				24	58049.91	57122.12	0.02	9016	14110	213597	368718
hanoi5	646	565	0.12	24	None	1880.90	NA	62844	73360	1109825	1324771
				24	None	1882.15	NA	517	522	1329967	1528334
				24	None	1880.52	NA	23770	27125	1146661	1384544
op.m2-z10-s4 ⁺	111	89	0.2	23	-33266.98	-33270.30	0.0	1446	12319	794	76102
				24	-33265.0	-33270.5	0.0	2520	2525	7491	61040
				23	-33269	-33269	0.0	2613	4631	5384	48433

Conclusions and future work

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- ▶ Can even be used to solve ILPs that are not PBOs, by treating the fractional part of each integer variable as a binary variable

Thank You!

Questions and comments to tamvadss@mcmaster.ca and
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